

Let  $f$  and  $g$  be functions and let  $C$  and  $D$  be propositions. By definition,

$$\begin{aligned} & y \\ = & \langle f : C \rangle \\ & z \\ = & \langle g : D \rangle \\ & x \end{aligned}$$

$$: C \wedge D \Rightarrow y = (f \cdot g).x$$

If we let

$$\begin{aligned} & y \\ = & \langle f : C \rangle \\ & x \end{aligned}$$

$$\begin{aligned} & y \\ = & \langle f : C \rangle \\ = & \langle id : \mathbf{true} \rangle \\ & x \end{aligned}$$

we can see that

$$\begin{aligned} & y \\ = & \langle f : C \rangle \\ & x \end{aligned}$$

$$\in C \Rightarrow y = f.x$$

We will use this proof format attain the following:

# Objective

The design of a demultiplexor chip together with a proof of its correctness.

## NTT5 Introduction

We wish to construct  $dmux$ , a so-called *demultiplexor*.

$dmux$  inputs a pair of booleans and outputs a pair of booleans i.e.

$$dmux : 2 \rightarrow 2$$

and is specified by

$$\llbracket dmux \rrbracket.[s, in] \triangleq [s \wedge in, \neg s \wedge in]$$

In order to construct  $dmux$ , we have the following chips at our disposal:

- ▶ nand  $\bar{A}B$ , not  $\neg C$ , and  $A \wedge B$ , or  $A \vee B$ , xor  $\neq C$
- ▶  $mux$ .

Hence we calculate:

$$\begin{aligned}
& [s \wedge in, \neg s \wedge in] \\
= & \langle \llbracket \text{second } \neg C \rrbracket : \llbracket \text{second } f \rrbracket . [a, b] = [a, \llbracket f \rrbracket . b] \rangle \\
& [s \wedge in, s \wedge in] \\
= & \langle \llbracket \text{fork} \bowtie f \rrbracket : \llbracket \text{fork} \bowtie f \rrbracket . [a] = [a, a] \rangle \\
& [s \wedge in] \\
= & \langle \llbracket \wedge C \rrbracket : \llbracket \wedge C \rrbracket . [a, b] = [a \wedge b] \rangle \\
& [s, in]
\end{aligned}$$

Recalling that

$$[ \llbracket f \rrbracket \gg g ] = \llbracket g \rrbracket \cdot \llbracket f \rrbracket ]$$

we arrive at our final program:

$$dmux \triangleq \wedge C \gg \text{fork} \bowtie \gg \text{second } \neg C$$